Ruzica Piskac, Ph.D. Leander Tentrup, M.Sc. Michael Gerke, M.Sc. Felix Klein, B.Sc.

Embedded Systems

Please indicate your **name**, **group number**, and **discussion slot tutor**. Only one submission per group is necessary.

Problem 1: Periodic Scheduling

For each of the following tasks sets, (1) determine whether an EDF-schedule and/or an RM-schedule exists, and (2) formally prove your answer.

$\Gamma = \{\tau_1, \tau_2, \tau_3\}$	$T_1 = D_1 = 3$ $T_2 = D_2 = 4$ $T_3 = D_3 = 8$	$C_1 = 1$ $C_2 = 2$ $C_3 = 1$
$\Delta = \{\tau_1, \tau_2, \tau_3\}$	$T_1 = D_1 = 2$ $T_2 = D_2 = 3$ $T_3 = D_3 = 4$	$C_1 = 1$ $C_2 = 1$ $C_3 = 1$
$\Pi = \{\tau_1, \tau_2, \tau_3, \tau_4\}$	$T_1 = D_1 = 2 T_2 = D_2 = 5 T_3 = D_3 = 8 T_4 = D_4 = 20$	$C_1 = 1$ $C_2 = 1$ $C_3 = 2$ $C_4 = 1$

Problem 2: Aperiodic Scheduling

Consider the following scheduling problem $1 \mid sync \mid T_w$:

Using a uniprocessor machine, find a schedule for a set $\mathcal{J} = \{J_1, \ldots, J_n\}$ of n synchronous tasks with computation times C_1, \ldots, C_n that minimizes the weighted sum of the completion times

$$T_w = \sum_{i=1}^n (w_i f_i) \,,$$

where $w_i > 0$ is a weight, and f_i is the time at which task *i* finishes its execution. (*Note:* The schedule is not required to respect the deadlines. We are only interested in minimizing T_{w} .)

- (a) Let \mathcal{J} be a task set, and let σ be a schedule for \mathcal{J} that is optimal with respect to the problem $1 \mid sync \mid T_w$. Formally prove that there exists a nonpreemptive schedule σ^* for \mathcal{J} with the same T_w of σ .
- (b) Devise a polynomial-time algorithm that, given a task set $\mathcal{J} = \{J_1, \ldots, J_n\}$, computes a schedule σ for \mathcal{J} that is optimal with respect to the scheduling problem $1 \mid sync \mid T_w$.
- (c) Formally prove that your algorithm computes an optimal schedule.

Problem 3: Optimality of Aperiodic Scheduling

Consider the problem of scheduling a set of synchronous tasks on a uniprocessor machine. It was shown in class that Jackson's EDD algorithm minimizes the maximum lateness

$$L_{max} = \max_{i} (f_i - d_i) \,.$$

For each of the following criteria, determine whether the EDD algorithm minimizes it:

- (a) average response time $\overline{R} = \frac{1}{n} \sum_{i=1}^{n} f_i$;
- (b) total completion time $T_c = \max_i(f_i)$;
- (c) weighted sum of completion times $T_w = \sum_{i=1}^n w_i f_i$;
- (d) number of late tasks $N_{late} = \sum_{i=1}^{n} (if \ d_i > f_i \ then \ 1 \ else \ 0).$

For each criterion, if EDD minimizes it, give a formal proof; otherwise give a counterexample.

Problem 4: Project Schedule

Make a schedule for the last project and present it in the discussion slot **this week** (10./11.07.). Indicate in this schedule when you plan to start with the subproblems and estimate a finishing time. Make sure to identify parallelizable tasks.