

# Automata, Games & Verification

Summary #5

# Complementation of Büchi Automata

## **Lemma 1.**

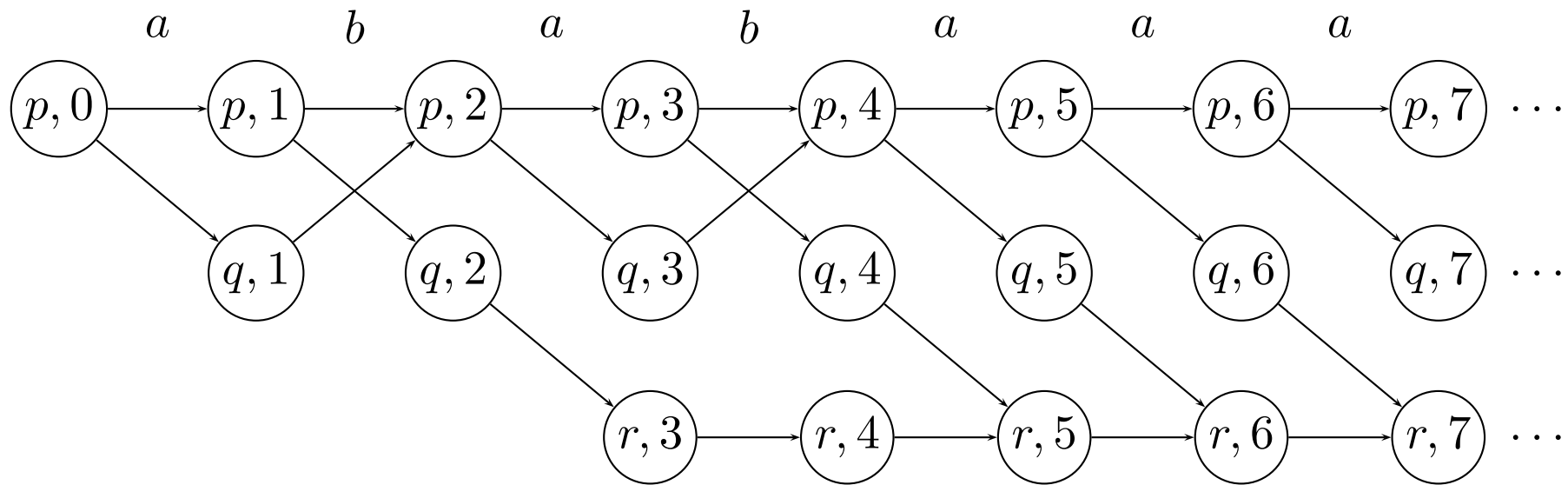
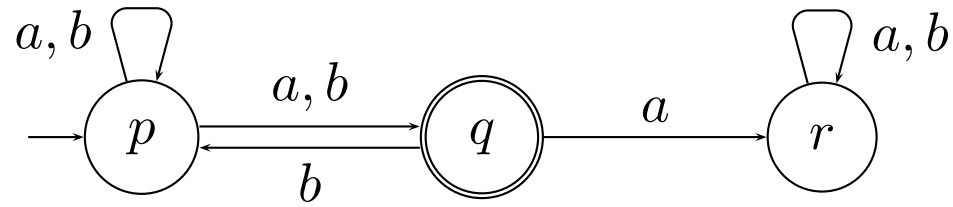
*If there exists an odd ranking for  $G$ , then  $\mathcal{A}$  does not accept  $\alpha$ .*

Let  $G'$  be a subgraph of  $G$ . We call a vertex  $\langle s, l \rangle$

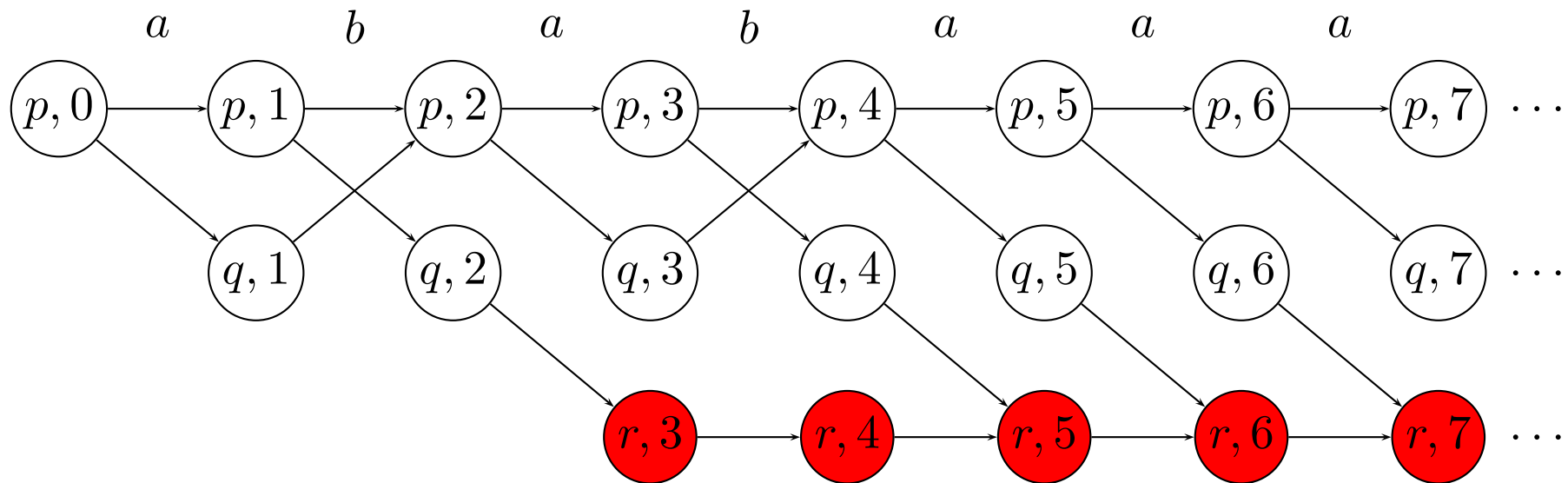
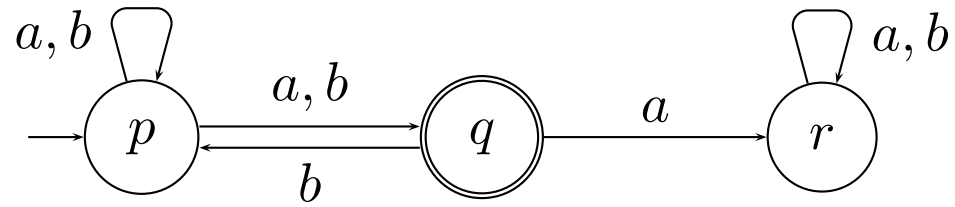
- **safe** in  $G'$  if for all vertices  $\langle s', l' \rangle$  reachable from  $\langle s, l \rangle$ ,  $s' \notin F$ , and
- **endangered** in  $G'$  if only finitely many vertices are reachable.

We define an infinite sequence  $G_0 \supseteq G_1 \supseteq G_2 \supseteq \dots$  of DAGs inductively as follows:

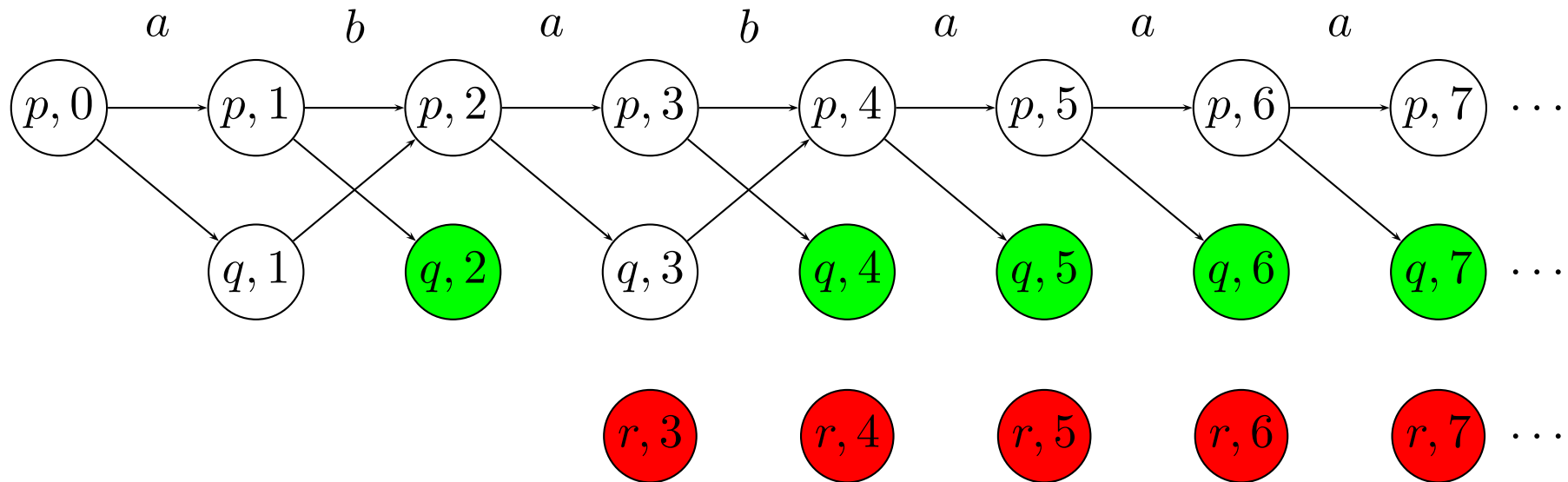
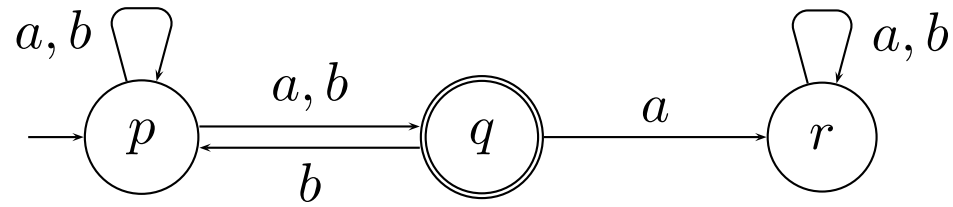
- $G_0 = G$
- $G_{2i+1} = G_{2i} \setminus \{ \langle s, l \rangle \mid \langle s, l \rangle \text{ is endangered in } G_{2i} \}$
- $G_{2i+2} = G_{2i+1} \setminus \{ \langle s, l \rangle \mid \langle s, l \rangle \text{ is safe in } G_{2i} \}$ .



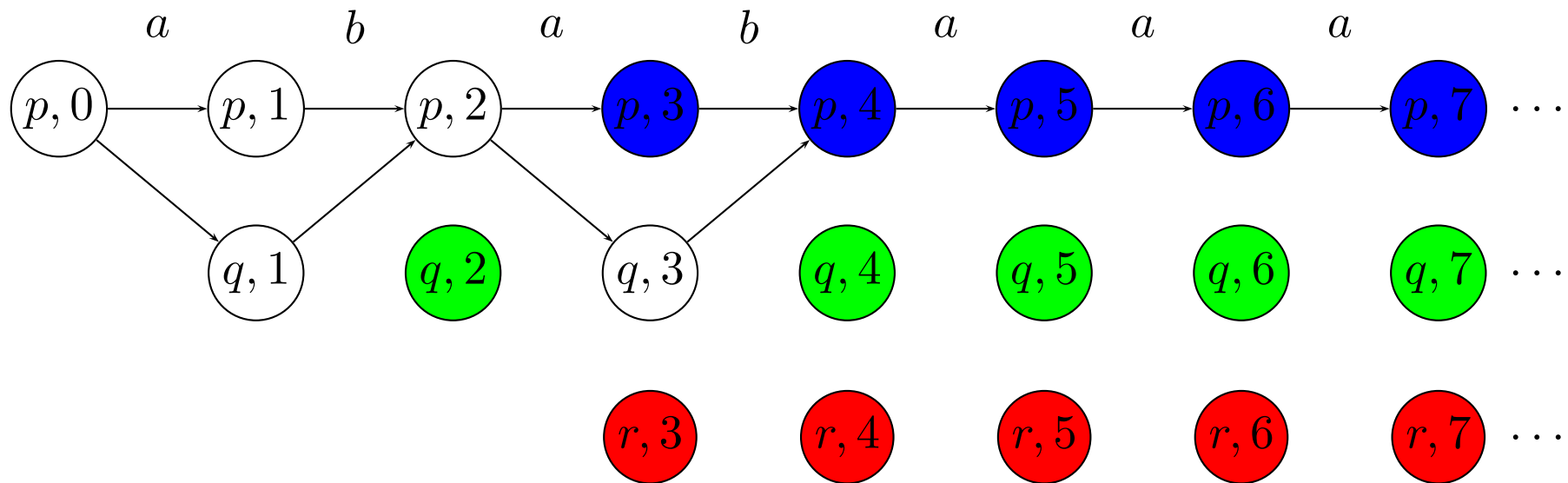
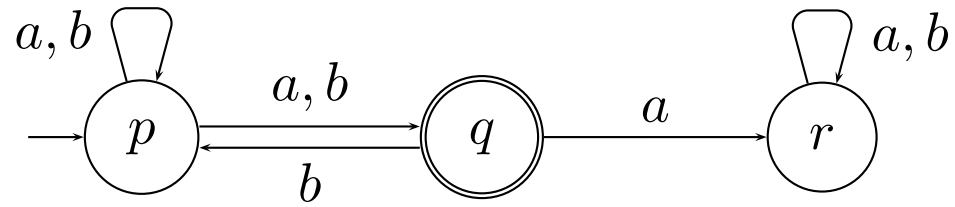
$$G = G_0 = G_1$$



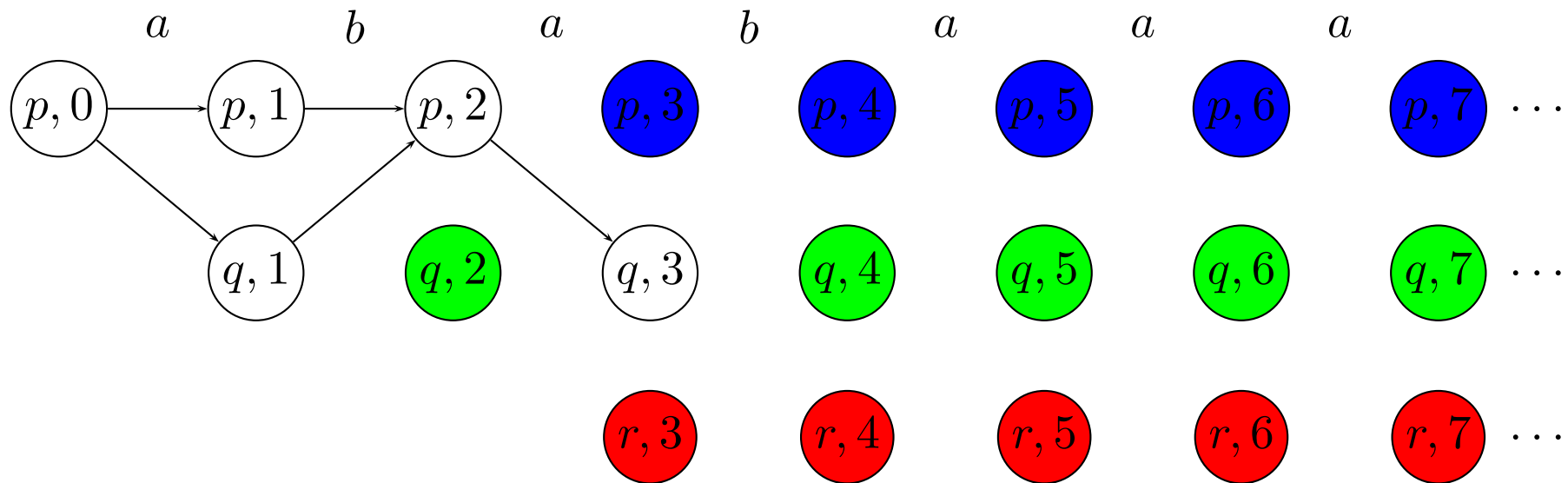
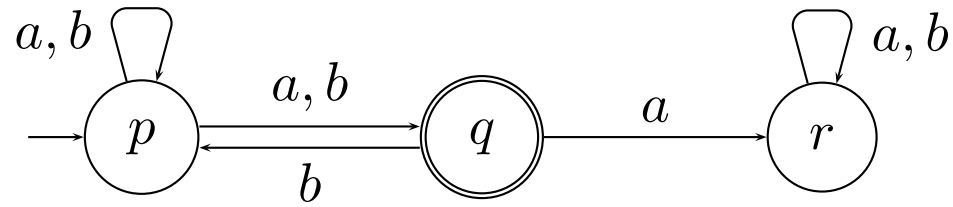
$G_1$



$G_2$

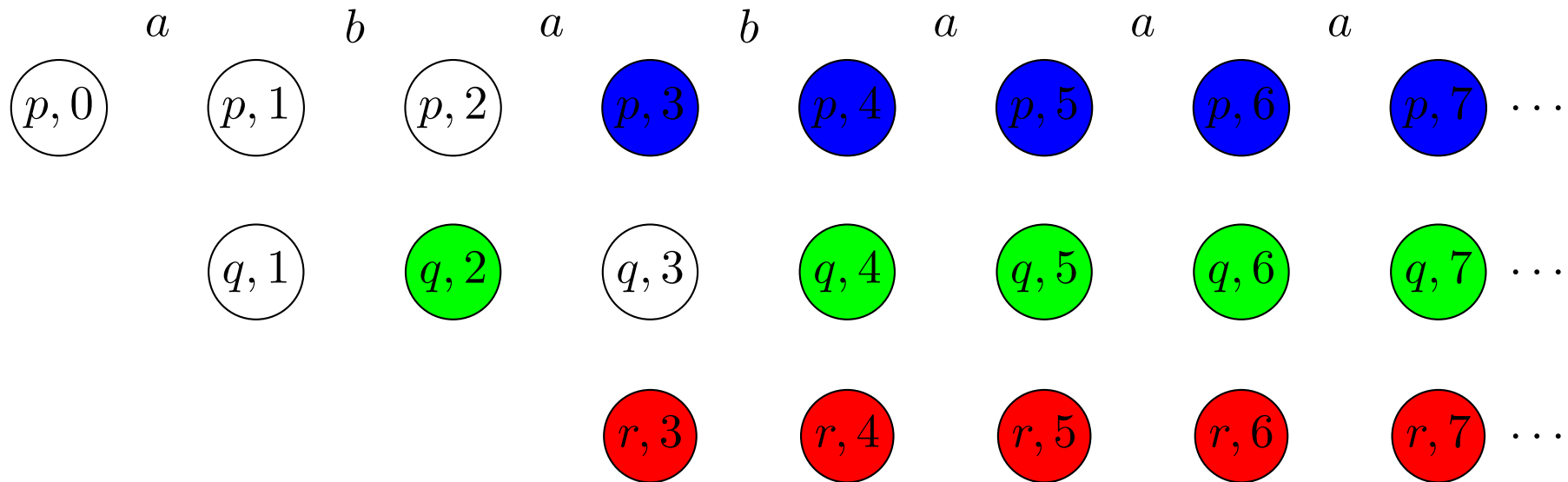
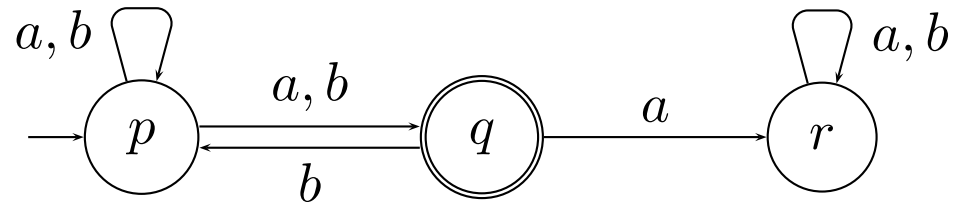


$G_3$



$G_4$





$G_5$

**Lemma 2.**

*If  $\mathcal{A}$  does not accept  $\alpha$ , then the following holds:*

*For every  $i \geq 0$  there exists an  $l_i$  such that*

*for all  $j \geq l_i$  at most  $|S| - i$  vertices of the form  $\langle -, j \rangle$  are in  $G_{2^i}$ .*

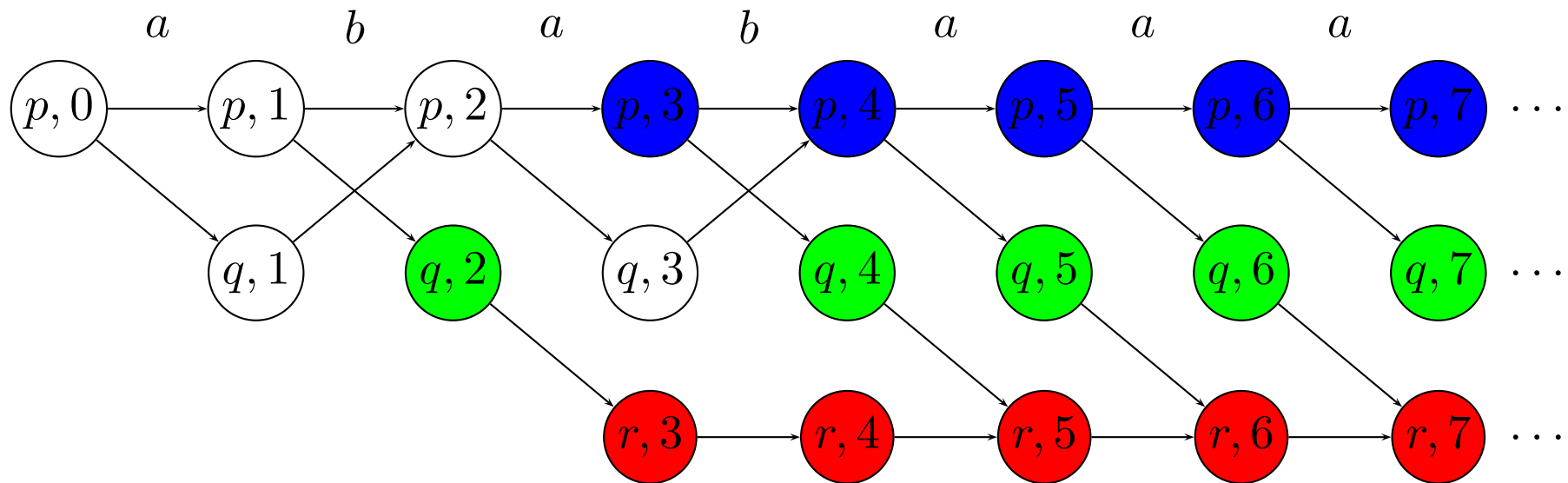
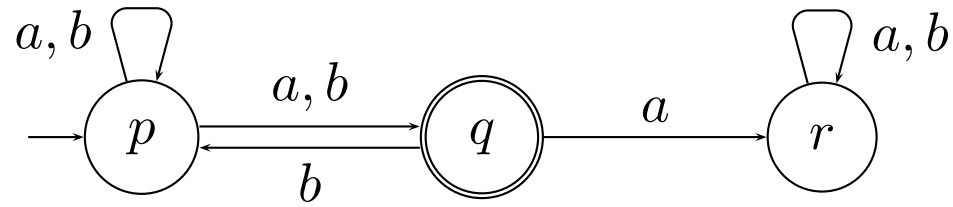
Proof by induction on  $i$ .

### **Lemma 3.**

*If  $\mathcal{A}$  does not accept  $\alpha$ , then there exists an odd ranking for  $G$ .*

We defined the ranking  $f$  where:

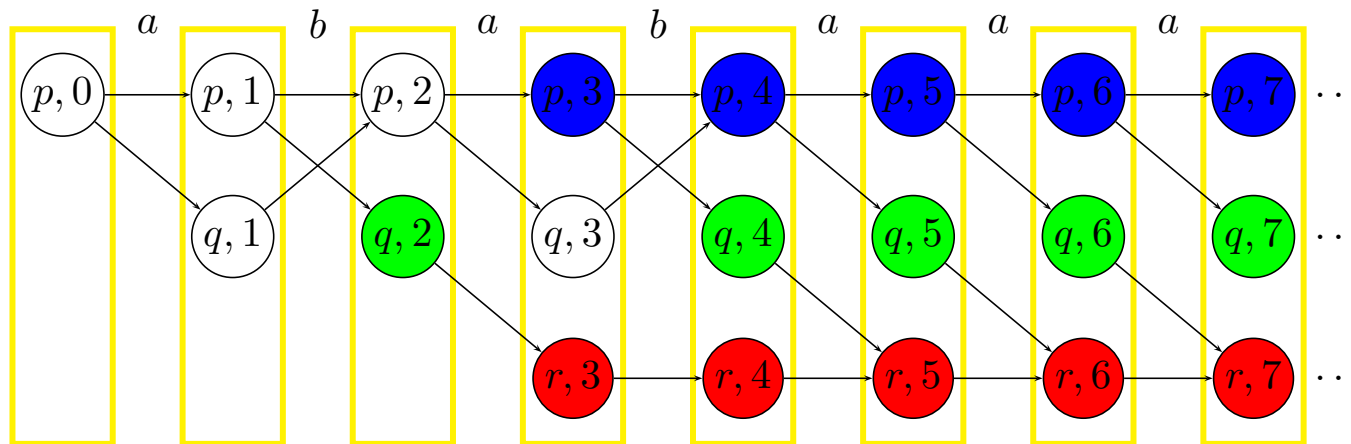
- $f(\langle s, l \rangle) = 2i$  if  $\langle s, l \rangle$  is endangered in  $G_{2i}$
- $f(\langle s, l \rangle) = 2i + 1$  if  $\langle s, l \rangle$  is safe in  $G_{2i+1}$



rank 1 — rank 2 — rank 3 — rank 4

# Complementation

**Theorem 1.** For each Büchi automaton  $\mathcal{A}$  there exists a Büchi automaton  $\mathcal{A}'$  such that  $\mathcal{L}(\mathcal{A}') = \Sigma^\omega \setminus \mathcal{L}(\mathcal{A})$ .



## Determinizing

**Theorem 2.** *The language  $(a \cup b)^*b^\omega$  is not recognizable by a deterministic Büchi automaton.*