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Embedded Systems 08/09 - Problem Set 9

Problem 1 - Abstract caches

An abstract cache can be represented as a partial function $f : \mathbb{N}_0 \rightarrow 2^V$ that maps age values (positive integers) to a subset of a finite object set *V*. Using this representation having the object set $V = \{a, b, c, d, e\}$, the abstract cache with 4 entries

$\{a\}$
$\{b,c\}$
$\{d\}$
$\{e\}$

can be represented as the partial function f with

$$f(0) = \{a\},\$$

$$f(1) = \{b, c\},\$$

$$f(2) = \{d\}, \text{ and }\$$

$$f(3) = \{e\}.$$

Assuming an LRU caching strategy, your task is to prove the following claims:

1. For an arbitrary abstract *must* cache f and an object $x \in V$, let f' be the resulting abstract *must* cache after accessing x in f. Then the following holds true:

$$\forall i. (x \in f(i) \implies f(i) \setminus \{x\} \subseteq f'(i))$$

(15 pts.)

2. For an arbitrary abstract *may* cache f with n entries and an object $x \in V$, let f' be the resulting abstract *may* cache after accessing x in f. Then the following holds true:

$$f'(0) = \{x\} \land \forall i < n-1. (f'(i+1) = f(i) \setminus \{x\})$$

(15 pts.)

(30 pts.)

Problem 2 - Cache analysis and predictability

(20 pts.)

FIFO is an alternative cache-replacement policy. In case of a cache-miss, the oldest cache entry is removed. Consider the following sequence of accesses:



- 1. Assume an LRU-cache with 2 entries. Apply a must-cache and a may-cache analysis of the program. Can you predict the exact cache state at the last access to *x*? (10 pts.)
- 2. Assume a FIFO-cache with 2 entries and an empty cache at the beginning. Can you predict a cache-hit or a cache-miss at the last access to *x*? (10 pts.)
- 3. Repeat (2) with an unknown cache-state at the beginning of the program. What can you say about the predictability of an LRU-cache compared to a FIFO-cache? (10 pts.)

Problem 3 - Abstract Interpretation

For a positive integer $n \in \mathbb{N}_0$ consider the lattices

$$C = (2^{\mathbb{N}_0}, \subseteq, \top_C, \bot_C) \text{ and} A = (2^{\{0,\dots,n-1\}}, \subseteq, \top_A, \bot_A),$$

as well as the two transfer functions

$$\begin{array}{rcl} \alpha_n & : & 2^{\mathbb{N}_0} \longrightarrow 2^{\{0,\dots,n-1\}} & \text{and} \\ \gamma_n & : & 2^{\{0,\dots,n-1\}} \longrightarrow 2^{\mathbb{N}_0} \end{array}$$

which are defined as

$$\begin{aligned} \alpha_n(c) &= \{ x' \in \mathbb{N}_0 \mid \exists x \in c. \ x' \equiv x \bmod n \}, \\ \gamma_n(a) &= \{ x' \in \mathbb{N}_0 \mid \exists x \in a \ \exists k \in \mathbb{N}_0. \ x' = kn + x \}, \end{aligned}$$

for any $c \in 2^{\mathbb{N}_0}$ and $a \in 2^{\{0,...,n-1\}}$.

- 1. Give suitable top and bottom elements \top_C , \perp_C , \top_A , and \perp_A . (5 pts.)
- 2. Prove that (α_n, γ_n) is a Galois connection. (15 pts.)

Problem 4 - Path analysis

Consider the following (pseudo code) program and its corresponding local WCET:

		MOTT
	Statement	WCEI
if (a) then	a	10
if (b) then	b	2
c ; d	С	4
elseif (e) then	d	3
f	е	1
else	f	8
g ; h	g	6
else	h	2
i ; j	i	7
	j	3

- 1. Give the control flow graph (in a graphical notation). (5 pts.)
- 2. Give an integer linear program for maximizing the overall WCET. (10 pts.)
- 3. Compute the overall WCET and its corresponding path. (5 pts.)