## Embedded Systems 08/09 - Problem Set 7

## Problem 1 - Priority Ceiling Protocol

Consider the Priority Ceiling Protocol. Using this protocol, give a picture describing a run of three tasks on one processor:

- Task 1 has the highest priority. Task 1 arrives at time $\mathrm{t}=5$. Task 1 consists of normal computation for 1 time unit, followed by critical section 1 for 1 time unit, followed by normal computation for 1 time unit.
- Task 2 has lower priority than Task 1. Task 2 arrives at time $t=2$. Task 2 consists of normal computation for 1 time unit, followed by critical section 2 for 2 time units, followed by normal computation for 1 time unit, followed by critical section 3 for 1 time unit, followed by normal computation for 1 time unit.
- Task 3 has the lowest priority. Task 3 arrives at time $t=0$. Task 3 consists of normal computation for 1 time unit, followed by critical section 3 for 2 time units, followed by normal computation for 1 time unit.

Your picture should depict which task is executed (and the type of computation: either normal or the critical section the task is in) in the processor at which point in time, covering the interval from time $t=0$ to time $t=13$. You should also explicitly give any changes in the priority of the Tasks.

## Problem 2-Scheduling

Consider the following scheduling problem $1 \mid$ sync $\mid \bar{R}$ :
Using a uniprocessor machine, find a schedule for $n$ aperiodic tasks $J_{1}, \ldots, J_{n}$ with computation times $C_{1}, \ldots, C_{n}$ that minimizes the average response time

$$
\bar{R}=\frac{1}{n} \sum_{i=1}^{n}\left(f_{i}-a_{i}\right),
$$

where $f_{i}$ is the time at which task $i$ finishes its execution. Assume that all tasks arrive at time 0 , that is, $a_{i}=0$, for all $i=1, \ldots, n$.
(a) Devise a scheduling algorithm that is optimal for the scheduling problem $1 \mid$ sync $\mid \bar{R}$. (20 pts.)
(b) Formally prove that your algorithm is optimal. Hint: If your algorithm is deterministic, then you can start by proving the following lemma: Let $S$ be a task-set, let $\sigma(S)$ be the schedule obtained by your algorithm, and let $\tau$ be any schedule. Then there exists a sequence of schedules $\tau_{0}, \ldots, \tau_{k}$ such that (i) $\tau_{0}=\tau$, (ii) $\tau_{k}=\sigma(S)$, and (iii) the average response time of $\tau_{i+1}$ is not greater than the average response of $\tau_{i}$, for $i=0, \ldots, k-1$. (30 pts.)

## Problem 3 - Multi-processor scheduling

Consider the following task sets:

- Task set 1 :

| Arrival time | Computation time | Deadline |
| :--- | :--- | :--- |
| 0 | 2 | 4 |
| 0 | 3 | 4 |
| 0 | 1 | 2 |
| 0 | 5 | 6 |

- Task set 2 :

| Arrival time | Computation time | Deadline |
| :--- | :--- | :--- |
| 0 | 3 | 5 |
| 0 | 4 | 6 |
| 0 | 2 | 7 |
| 0 | 5 | 8 |

Please answer the following questions:

1. For each of the task sets, either give a feasible schedule for two processors or prove that there exists no feasible schedule. ( $10+10$ pts.)
2. For each task set, also answer the question whether the result (i. e. whether the task set is schedulable or not schedulable) you have obtained in part one of this problem could be different if the arrival times would be different. ( $5+5 \mathrm{pts}$.)
